## Lecture 13: Information Theory and Counting: Shearer's Lemma

- Entropy = "Surprise"
- $X$ is a random variable
- $p_{X}(x)$ is the probability of $x$ according to the distribution $X$
- Surprise of seeing $x$ is, $h_{X}(x):=-\log p_{X}(x)$
- Entropy of $X$ is the expected surprise, $H(X):=\mathbb{E}_{X \sim X}\left[h_{X}(x)\right]$
- Alternately, $H(X)=\sum_{x}-p_{X}(x) \log p_{X}(x)$


## Conditional Entropy

- $(X, Y)$ is a joint distribution
- $(X \mid Y=y)$ is the conditional distribution of $X$ conditioned on the fact that $Y=y$
- Entropy of $(X \mid Y=y)$ is defined by $H(X \mid Y=y)$
- Conditional entropy $H(X \mid Y):=\mathbb{E}_{y \sim Y}[H(X \mid Y=y)]$


## Entropy

- Chain Rule: $H(X Y)=H(X)+H(X \mid Y)$
- Conditional Chain Rule: $H(X Y \mid Z)=H(X \mid Z)+H(Y \mid X Z)$
- Inequalities:
- $0 \leqslant H(X) \leqslant|r a n g e(X)|$
- $H(X) \geqslant H(X \mid f(Y)) \geqslant H(X \mid Y)$
- Binary Entropy Function:
$h(p):=-p \log p-(1-p) \log (1-p)$


## Binomial Coefficient Tail

## Theorem

$$
\sum_{i \leqslant \alpha n}\binom{n}{i} \leqslant 2^{h(\alpha) n}
$$

- Let $\mathcal{C}$ be the set of all subsets of $[n]$ of size at most $\alpha n$
- Let $X$ be a uniform distribution over $\mathcal{C}$
- Let $\left(X_{1}, \ldots, X_{n}\right)$ be the characteristic vector corresponding to the subset sampled by $X$
- $\log |\mathcal{C}|=H(X)=H\left(X_{1}, \ldots, X_{n}\right)=\sum_{i \in[n]} H\left(X_{i} \mid X_{<i}\right) \leqslant$ $\sum_{i \in[n]} H\left(X_{i}\right)$
- Since all indices are symmetric, $H\left(X_{i}\right)=H\left(X_{1}\right)$
- Note that $H\left(X_{1}| | X \mid=i\right)=h(i / n) \leqslant h(\alpha)$, for $i \leqslant \alpha n$
- Therefore, $H\left(X_{1}\right) \leqslant h(\alpha)$
- Overall $\log |\mathcal{C}| \leqslant n h(\alpha)$


## Identifying Bad Balls

Consider the task of designing a set $\mathcal{D}=\left\{D_{1}, \ldots, D_{\ell}\right\}$ such that $D_{i} \subseteq[n]$, for $i \in[\ell]$ such that:

- Consider $n$ ordered balls
- Let $B$ be the set of positions with bad balls
- Suppose we are given an oracle that on input $D \subseteq[n]$ outputs $|B \cap D|$
- Using each set in $\mathcal{D}$ to query the oracle, output $B$


## Theorem

$$
\ell \geqslant n / \log (n+1)
$$

- Note that $B \mapsto\left(\left|B \cap D_{1}\right|, \ldots,\left|B \cap D_{\ell}\right|\right)$ is a bijection (two different $B$ and $B^{\prime}$ cannot have the same sequence, otherwise we cannot distinguish $B$ from $B^{\prime}$ )
- Let $X$ be a uniform random variable over $2^{[n]}$ (i.e., the set of all subsets of $[n]$ )
- $n=H(X) \leqslant \sum_{i \in[\ell]} H\left(\left|X \cap D_{i}\right|\right) \leqslant \ell \log (n+1)$


## Number of Matchings

Let $G=(A, B, E)$ be a bipartite graph

## Theorem (Brégman's Theorem)

Number of perfect matchings in $G$ is at most $\prod_{v \in A}(d(v)!)^{1 / d(v)}$

- Let $\Sigma$ be the set of all perfect matchings
- Let $\sigma$ be a uniform random variable over $\Sigma$
- $\log |\Sigma|=H(\sigma)=\sum_{v \in A} H\left(\sigma(v) \mid \sigma(u)_{u<v}\right)$
- Trivial upper bound by $H(\sigma(v)) \leqslant \log d(v)$
- Idea: Expand according to a random permutation $\tau$ of vertices in $A$
- Think: How to get $\frac{1}{d(v)}(1+\cdots+d(v))$ as upper bound to get the result
- Consider a set $S$ of $n$ points in 3-dimensions
- Let $n_{1}$ be number of unique points by projecting $S$ on $X=0$ plane, $n_{2}$ be the number of unique points by projecting $S$ on $Y=0$ plane and $n_{3}$ be the number of unique points by projecting $S$ on $Z=0$ plane


## Theorem

$$
n \leqslant\left(n_{1} n_{2} n_{3}\right)^{1 / 2}
$$

- Let $(X, Y, Z)$ represent the coordinates of uniformly chosen point in $S$
- $\log n=H(X, Y, Z)=H(X)+H(Y \mid X)+H(Z \mid X Y)$
- $\log n_{1} \geqslant H(Y, Z)=H(Y)+H(Z \mid Y) \geqslant H(Y \mid X)+H(Z \mid X Y)$
- $\log n_{2} \geqslant H(X, Z)=H(X)+H(Z \mid X) \geqslant H(X)+H(Z \mid X Y)$
- $\log n_{3} \geqslant H(X, Y)=H(X)+H(X \mid Y)$
- $\log n \leqslant \frac{1}{2} \log \left(n_{1} n_{2} n_{3}\right)$


## Shearer's Lemma

- Let $\mathcal{F}$ be a set of subsets of $[n]$
- For every $i \in[n]$, there are at least $t$ subsets in $\mathcal{F}$ that contain i


## Theorem (Shearer's Lemma)

$$
H\left(X_{1}, \ldots, X_{n}\right) \leqslant \frac{1}{t} \sum_{F \in \mathcal{F}} H\left(X_{F}\right)
$$

- "Sub-additivity of Entropy" is obtained by considering $\mathcal{F}$ as the set of all singleton sets
- "Volume computed by projections" is obtained by considering $\mathcal{F}$ as the set of all subsets of size $(n-1)$


## Theorem (Loomis-Whitney Theorem)

Let $B$ be a measurable body in $\mathbb{R}^{d}$ and $|\cdot|$ represent the volume. Let $B_{j}$, for $j \in[d]$, represent the body when $B$ is projected along the $j$-th coordinate axis. Then:

- The $i$-th smallest index in $F$ is represented by $F_{i}$
- Consider the manipulation similar to the "volume argument"

$$
\begin{aligned}
\sum_{F \in \mathcal{F}} H\left(X_{F}\right) & =\sum_{F \in \mathcal{F}} \sum_{i \leqslant|F|} H\left(X_{F_{i}} \mid X_{\left\{j<F_{i}\right\} \cap F}\right) \\
& \geqslant \sum_{F \in \mathcal{F}} \sum_{i \leqslant|F|} H\left(X_{F_{i}} \mid X_{\left\{j<F_{i}\right\}}\right) \\
& \geqslant t \sum_{i \in[n]} H\left(X_{i} \mid X_{<i}\right)=t H\left(X_{1}, \ldots, X_{n}\right)
\end{aligned}
$$

