# Lecture 13: Information Theory and Counting: Shearer's Lemma

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- Entropy = "Surprise"
- X is a random variable
- $p_X(x)$  is the probability of x according to the distribution X
- Surprise of seeing x is,  $h_X(x) := -logp_X(x)$
- Entropy of X is the expected surprise,  $H(X) := \mathbb{E}_{x \sim X}[h_X(x)]$
- Alternately,  $H(X) = \sum_{x} -p_X(x) \log p_X(x)$

- (X, Y) is a joint distribution
- (X|Y = y) is the conditional distribution of X conditioned on the fact that Y = y
- Entropy of (X|Y = y) is defined by H(X|Y = y)
- Conditional entropy  $H(X|Y) := \mathbb{E}_{y \sim Y}[H(X|Y = y)]$

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- Chain Rule: H(XY) = H(X) + H(X|Y)
- Conditional Chain Rule: H(XY|Z) = H(X|Z) + H(Y|XZ)
- Inequalities:

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$$0 \leq H(X) \leq |\operatorname{range}(X)|$$

- $H(X) \ge H(X|f(Y)) \ge H(X|Y)$
- Binary Entropy Function:

$$h(p) := -p \log p - (1-p) \log(1-p)$$

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## Binomial Coefficient Tail

#### Theorem

$$\sum_{\leqslant \alpha n} \binom{n}{i} \leqslant 2^{h(\alpha)n}$$

- Let C be the set of all subsets of [n] of size at most  $\alpha n$
- Let X be a uniform distribution over  $\mathcal C$
- Let  $(X_1, \ldots, X_n)$  be the characteristic vector corresponding to the subset sampled by X
- $\log |\mathcal{C}| = H(X) = H(X_1, \dots, X_n) = \sum_{i \in [n]} H(X_i | X_{< i}) \leq \sum_{i \in [n]} H(X_i)$
- Since all indices are symmetric,  $H(X_i) = H(X_1)$
- Note that  $H(X_1||X|=i) = h(i/n) \leqslant h(\alpha)$ , for  $i \leqslant \alpha n$
- Therefore,  $H(X_1) \leq h(\alpha)$
- Overall  $\log |\mathcal{C}| \leq nh(\alpha)$

# Identifying Bad Balls

Consider the task of designing a set  $\mathcal{D} = \{D_1, \ldots, D_\ell\}$  such that  $D_i \subseteq [n]$ , for  $i \in [\ell]$  such that:

- Consider *n* ordered balls
- Let *B* be the set of positions with bad balls
- Suppose we are given an oracle that on input  $D \subseteq [n]$  outputs  $|B \cap D|$
- $\bullet\,$  Using each set in  ${\cal D}$  to query the oracle, output B

### Theorem

$$\ell \geqslant n/\log(n+1)$$

- Note that B → (|B ∩ D<sub>1</sub>|,..., |B ∩ D<sub>ℓ</sub>|) is a bijection (two different B and B' cannot have the same sequence, otherwise we cannot distinguish B from B')
- Let X be a uniform random variable over 2<sup>[n]</sup> (i.e., the set of all subsets of [n])

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$$n = H(X) \leqslant \sum_{i \in [\ell]} H(|X \cap D_i|) \leqslant \ell \log(n+1)$$

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## Number of Matchings

Let G = (A, B, E) be a bipartite graph

Theorem (Brégman's Theorem)

Number of perfect matchings in G is at most  $\prod_{v \in A} (d(v)!)^{1/d(v)}$ 

- $\bullet\,$  Let  $\Sigma$  be the set of all perfect matchings
- Let  $\sigma$  be a uniform random variable over  $\Sigma$
- $\log |\Sigma| = H(\sigma) = \sum_{v \in A} H(\sigma(v) | \sigma(u)_{u < v})$
- Trivial upper bound by  $H(\sigma(v)) \leq \log d(v)$
- Idea: Expand according to a random permutation au of vertices in A
- Think: How to get  $\frac{1}{d(v)}(1+\cdots+d(v))$  as upper bound to get the result

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## Introduction to Shearer's Lemma

- Consider a set S of n points in 3-dimensions
- Let n<sub>1</sub> be number of unique points by projecting S on X = 0 plane, n<sub>2</sub> be the number of unique points by projecting S on Y = 0 plane and n<sub>3</sub> be the number of unique points by projecting S on Z = 0 plane

#### Theorem

 $n \leqslant \left(n_1 n_2 n_3\right)^{1/2}$ 

- Let (X, Y, Z) represent the coordinates of uniformly chosen point in S
- $\log n = H(X, Y, Z) = H(X) + H(Y|X) + H(Z|XY)$
- $\log n_1 \ge H(Y,Z) = H(Y) + H(Z|Y) \ge H(Y|X) + H(Z|XY)$
- $\log n_2 \ge H(X,Z) = H(X) + H(Z|X) \ge H(X) + H(Z|XY)$
- $\log n_3 \ge H(X, Y) = H(X) + H(X|Y)$
- $\log n \leq \frac{1}{2} \log(n_1 n_2 n_3)$

## Shearer's Lemma

Let F be a set of subsets of [n]
For every i ∈ [n], there are at least t subsets in F that contain i

### Theorem (Shearer's Lemma)

$$H(X_1,\ldots,X_n) \leqslant \frac{1}{t} \sum_{F \in \mathcal{F}} H(X_F)$$

- $\bullet$  "Sub-additivity of Entropy" is obtained by considering  ${\cal F}$  as the set of all singleton sets
- "Volume computed by projections" is obtained by considering  $\mathcal{F}$  as the set of all subsets of size (n-1)

### Theorem (Loomis-Whitney Theorem)

Let B be a measurable body in  $\mathbb{R}^d$  and  $|\cdot|$  represent the volume. Let  $B_j$ , for  $j \in [d]$ , represent the body when B is projected along the j-th coordinate axis. Then:

- The *i*-th smallest index in F is represented by  $F_i$
- Consider the manipulation similar to the "volume argument"

$$\sum_{F \in \mathcal{F}} H(X_F) = \sum_{F \in \mathcal{F}} \sum_{i \leq |F|} H(X_{F_i} | X_{\{j < F_i\} \cap F})$$
  
$$\geq \sum_{F \in \mathcal{F}} \sum_{i \leq |F|} H(X_{F_i} | X_{\{j < F_i\}})$$
  
$$\geq t \sum_{i \in [n]} H(X_i | X_{< i}) = tH(X_1, \dots, X_n)$$

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